

Attainable Knowledge and Omniscience

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The paper investigates an evidence-based semantics for epistemic logics. It is shown that the properties of knowledge obtained from a potentially infinite body of evidence are described by modal logic S5. At the same time, the properties of knowledge obtained from only a finite subset of this body are described by modal logic S4. The main technical result is a sound and complete bi-modal logical system that describes properties of these two modalities and their interplay.

1 Introduction

The distinction between *potential* and *actual* infinity goes back to Zeno's paradoxes. The former refers to constructions involving a potentially unlimited, but finite, number of steps while the latter refers to an infinite set as a whole.

There has been a long tradition of distinguishing these two types of infinity in logic. We say that a set is decidable if the membership in the set can be verified through a potentially unlimited but finite number of steps of a Turing machine. Most people believe that Peano Arithmetic is consistent and, thus, its consistency could be verified by observing that none of the countably many finite-length derivations ends with a contradiction. Yet, Gödel's second incompleteness theorem claims that the consistency cannot be established by a single finite-length derivation or a finite set of such derivations. More generally, logicians usually accept admissible inference rules with finitely many assumptions and reject ω -rules that deduce statements from infinitely many assumptions. A compactness property for a logical system states that if a statement semantically follows from an infinite set of assumptions, then it semantically follows from a potentially unlimited but finite subset of the assumptions.

In this paper we investigate the difference between the knowledge that can be derived from a potentially unlimited, but finite, subset of the body of evidence and the knowledge that can be inferred from the whole infinite body of evidence. We refer to the first form of knowledge as *attainable knowledge* and to the second as *omniscience*. We show that the properties of the attainable knowledge are captured by modal logic S4, while those of the omniscience are given by modal logic S5.

1.1 Epistemic Logics S4 and S5

Epistemic logic S5 is an extension of logic S4 by the Negative Introspection axiom: $\neg K\phi \rightarrow K\neg K\phi$, where K is the knowledge modality. Informally, this axiom expresses an agent awareness of the limits of her knowledge: if an agent does not know something, then she must know that she does not know it. Even without taking into account the size of the body of supporting evidence, the validity of this axiom has long been a subject for philosophical discussions. Hintikka [8] rejects this axiom. Stalnaker [15] cites a belief-based argument against it. The argument relies on three standard assumptions about knowledge and beliefs: (i) beliefs of an agent are consistent, (ii) an agent believes in anything she knows, and (iii) anything an agent knows is true. We interpret Stalnaker's argument in terms of a mathematician's belief

about her knowledge. Imagine a mathematician who spent last several years working on a conjecture φ . One day she announces “I believe I now know that the conjecture is true”¹. We write this as $BK\varphi$. Because her beliefs are consistent, we must conclude that she does not believe in the opposite: $\neg B\neg K\varphi$. Hence, $\neg K\neg K\varphi$ by the contraposition of assumption (ii). Thus, by the contrapositive of the Negative Introspection axiom, we have $K\varphi$. Then, φ is true by assumption (iii). Following Stalnaker, we have just shown that if a mathematician believes that she knows that conjecture φ is true, then conjecture φ must be true. However, intuitively, this cannot be true because her belief alone should not make the conjecture true.

While the Negative Introspection axiom might not be true as the most general epistemic axiom, Fagin, Halpern, Moses, and Vardi point out that it is “appropriate for many multi-agent systems applications”, such as the analysis of the muddy children puzzle, knowledge bases, games of imperfect information, and message-passing systems [6, p.18]. In spite of its limitations, S5 has been accepted as the standard epistemic logic in the modern literature [1, 2, 5, 11]. There also have been suggestions to consider modal systems in between S4 and S5 [9], [10, p.82].

The purpose of this paper is not to settle the discussion about the merits of S4 vs. S5, but to highlight the difference between these two systems from the viewpoint of an evidence-based semantics. The main technical result of our work is a sound and complete logical system with two modalities that capture individual properties of the attainable knowledge and of the omniscience as well as the properties that connect these two types of knowledge.

1.2 Grand Hotel Example

Consider the famous Hilbert Grand Hotel that has infinitely many rooms. By a single observation in this setting we mean examining a room in order to establish whether it is empty. If the hotel has vacancies, then at least one room is empty. By opening just a single door of that specific room one can learn that the hotel has vacancy. We write A (“the hotel has vacancy”) to express this fact. The modality A denotes *attainable knowledge* that can be formed from finitely many observations. In fact, in this case a single observation suffices. On the other hand, if the hotel has no vacancy, then the knowledge about this cannot be obtained from examining any finite subset of the rooms: $\neg A$ (“the hotel has no vacancy”). At the same time, one can learn that the hotel is full by examining all rooms in the hotel. We write this as O (“the hotel has no vacancy”), where modality O denotes the *omniscience* that can be formed from a potentially infinite body of evidence.

Let us first show that attainable knowledge does not satisfy the Negative Introspection axiom:

$$\neg A\varphi \rightarrow A\neg A\varphi. \quad (1)$$

Consider the epistemic world w in which the hotel is full and let φ be statement “the hotel has vacancy”. Thus, $w \not\models \varphi$. Hence, one cannot know that φ is true from examining (finitely many) hotel rooms: $w \not\models A\varphi$. Hence, $w \models \neg A\varphi$. At the same time, informally, the only reason why $w \models \neg A\varphi$ is true is because the hotel is full. Since the latter cannot be established through finitely many observations, $w \models \neg A\neg A\varphi$. More formally, to prove $w \models \neg A\neg A\varphi$ we need to show that no matter which finite set of evidence is examined, there still will be an epistemic world u indistinguishable from w such that $u \models A\varphi$. Indeed, suppose that we have chosen to examine rooms r_1, \dots, r_n . Let r be any room different from rooms r_1, \dots, r_n . Let u be an epistemic world in which all rooms except for room r are occupied. Thus, epistemic worlds w and u are indistinguishable through the chosen finite set of examinations. Yet, $u \models A\varphi$

¹In Stalnaker’s words, “agent ... take[s] herself to know” that the conjecture is true [15, p.177].

because in epistemic world u , it is enough to examine a single room r to learn that the hotel has vacancy. This concludes the counterexample for the Negative Introspection axiom (1).

From the classical (non-evidence-based) Kripke semantics point of view, the Negative Introspection axiom captures the fact that the indistinguishability relation on epistemic states is symmetric. In our evidence-based semantics, the indistinguishability by a single evidence is also symmetric. That is, if one can distinguish one state of the hotel from another state by examining just a single room, then one can also distinguish the latter state from the former by examining the same room. The reason why the Negative Introspection axiom fails, under the evidence-based semantics, is because *the indistinguishability relation between a world and a set of worlds by a potentially unlimited, but finite body of evidence is not symmetric*. Indeed, consider the epistemic world in which the hotel is empty and the set of all worlds in which the hotel is not empty. By examining a right single room in a non-empty hotel, one can distinguish it from an empty hotel. Yet, after examining any finite subset of rooms in an empty hotel one cannot distinguish the current world from the set of worlds in which the hotel is not empty.

Let us now consider a weaker form of the Negative Introspection axiom used in logic S4.4 [9]:

$$\varphi \rightarrow (\neg A\varphi \rightarrow A\neg A\varphi). \quad (2)$$

The above counterexample for formula (1) does not work as a counterexample for formula (2) because in that setting $w \not\models \varphi$. Nevertheless, principle (2) is not valid in general. Indeed, let us modify the hotel setting by assuming that some rooms in the hotel might have bedbugs. The presence of the bedbugs can be tested when a room is examined. Furthermore, let us assume that once a single room in the hotel becomes infected with bedbugs all guests immediately leave the hotel. Thus, the hotel might have either (a) visitors and no bedbugs, or (b) no visitors and no bedbugs, or (c) bedbugs and no visitors. Consider an epistemic world w in which the hotel has no visitors and no bugs. Let φ be the statement “the hotel has no visitors”. Thus, $w \models \varphi$. Note that there are two ways to verify that the hotel has no visitors: either by examining all rooms to observe that they are all vacant or by examining a single room that contains bedbugs. Since in the epistemic world w the hotel is not infected with bugs, the only way to verify that the hotel is empty in this world is to examine all rooms. Thus, $w \models \neg A\varphi$. To finish the counterexample for formula (2), it suffices to show that $w \not\models A\neg A\varphi$. In other words, we need to prove that, after examining a finite set of rooms r_1, \dots, r_n , we will not be able to distinguish epistemic state w from such an epistemic state u that $u \models A\varphi$. Indeed, let r be any room different from rooms r_1, \dots, r_n and let u be the epistemic world in which room r is infected with bedbugs. Note that $u \models A\varphi$ because in epistemic state u it is enough to examine the bug-infected room r to conclude that the hotel is empty.

1.3 Formal Semantics of Evidence

To make our Grand Hotel example more formal, we introduce the formal semantics of evidence-based knowledge. Since we view evidence as a way to distinguish epistemic worlds, in this paper we interpret the pieces of evidence as equivalence relations on the worlds. When an agent takes into account several pieces of evidence, equivalence relations corresponding to these pieces intersect to form the equivalence relation of the agent. In the second Grand Hotel example above, each room is in one of the three states: occupied, vacant without bedbugs, and vacant with bedbugs. An epistemic world can be described by specifying the state of each room. The observation that consists of examining room r can distinguish two epistemic worlds in which room r is in different states. It cannot distinguish two epistemic worlds in which room r is in the same state.

In other words, we assemble an agent’s knowledge from several pieces of evidence in a similar way as how distributed knowledge [6] of a group is assembled from the individual observations of the members

of the group. Thus, the logical system developed in this paper could also be used to describe two different forms of group knowledge by an infinite group of agents: modality O represents the standard distributed knowledge by the whole group and modality A represents the distributed knowledge by a finite subgroup of the whole group.

The formal evidence-based semantics described above is similar to the one for the budget-constrained knowledge proposed by us earlier [12, 13]. It is different from the neighbourhood semantics of evidence investigated by van Benthem and Pacuit [4] and the probabilistic approach of Halpern and Pucella [7]. It is not clear how and if the results in the current paper could be applied to these alternative semantics.

1.4 Logical System

In this paper we propose a sound and complete logical system that describes the universal properties of omniscience modality O and attainable knowledge modality A . The axioms involving only modality O are exactly those forming modal logic S5:

1. Truth: $O\varphi \rightarrow \varphi$,
2. Positive Introspection: $O\varphi \rightarrow OO\varphi$,
3. Negative Introspection: $\neg O\varphi \rightarrow O\neg O\varphi$,
4. Distributivity: $O(\varphi \rightarrow \psi) \rightarrow (O\varphi \rightarrow O\psi)$.

Later we do not list the Positive Introspection axiom among axioms of our system because, just like in the case of logic S5, the Positive Introspection axiom is derivable from the other axioms. We prove this in Lemma 2. The axioms involving only modality A are exactly those forming modal logic S4:

1. Truth: $A\varphi \rightarrow \varphi$,
2. Positive Introspection: $A\varphi \rightarrow AA\varphi$,
3. Distributivity: $A(\varphi \rightarrow \psi) \rightarrow (A\varphi \rightarrow A\psi)$.

Finally, there appears to be two independent axioms that capture the interplay of the two modalities:

1. Monotonicity: $A\varphi \rightarrow O\varphi$,
2. Mixed Negative Introspection: $\neg A\varphi \rightarrow O\neg A\varphi$.

The first of these axioms states that any knowledge that can be formed from a finite subset of observations can also be formed on the bases of the whole set of observations. The second axiom is significantly more interesting. We have seen in our first Grand Hotel example that the Negative Introspection axiom is not true for attainable knowledge. Namely an agent in a fully occupied hotel cannot learn that statement A (“the hotel has vacancies”) is false by examining only finitely many rooms. It can learn this, however, by examining all room in the hotel and this is exactly what the Mixed Negative Introspection axiom claims. Surprisingly, the Mixed Negative Introspection axiom is provable from other axioms of our logical system. We show this in Lemma 3. As a result, Mixed Negative Introspection is not included as an axiom of our system. Additionally, it is easy to see that the Truth axiom for modality A follows from the Truth axiom for modality O and the Monotonicity axiom. For this reason, we do not list the Truth axiom for modality A as one of our axioms either. Finally, although one can state two forms of the Necessitation inference rule: one for modality O and another for modality A , the former follows from the latter and the Monotonicity axiom. Thus, our system, in addition to Modus Ponens, only includes the Necessitation rule for modality A .

1.5 Outline

The rest of the paper is organized as following. Section 2 defines the syntax and an evidence-based semantics of our logical system. Section 3 lists the axioms and the inference rules of this system and proves the Mixed Negative Introspection axiom mentioned above. In Section 4, we prove the soundness of our logical system with respect to the evidence-based semantics. In Section 5, we state the completeness theorem and define the canonical model used in its proof. The full proof of the completeness can be found in the appendix. Section 7 concludes.

2 Syntax and Semantics

In this section, we describe the formal syntax and the formal evidence-based semantics of our logical system. Throughout the rest of the paper we assume a fixed infinite set of propositional variables.

Definition 1 *The set Φ of all formulae φ is defined by grammar $\varphi := p \mid \neg\varphi \mid \varphi \rightarrow \varphi \mid A\varphi \mid O\varphi$, where p represents propositional variables.*

Definition 2 *An evidence model is $\langle W, E, \{\sim_e\}_{e \in E}, \pi \rangle$, where*

1. W is a (possibly empty) set of “epistemic worlds”,
2. E is an arbitrary (possibly empty) “evidence” set,
3. \sim_e is an “indistinguishability” equivalence relation on W for each $e \in E$,
4. π is a function that maps propositional variables into subsets of W .

For instance, in our first Grand Hotel example, an epistemic world is a function $\mathbb{N} \rightarrow \{\text{vacant}, \text{occupied}\}$ that assigns a state to each room in the hotel. The set of evidence is \mathbb{N} , where evidence with number $r \in \mathbb{N}$ corresponds to examining room with number r in this hotel. Epistemic worlds w_1 and w_2 are \sim_r -equivalent if room number r has the same state in both of the worlds. In other words, $w_1 \sim_r w_2$ if $w_1(r) = w_2(r)$. Finally a function π may, for example, map propositional variable p into the set of all epistemic worlds representing a nonempty hotel: $\pi(p) = \{w \in W \mid \exists r \in \mathbb{N} \text{ such that } w(r) = \text{occupied}\}$.

In this paper, we write $w_1 \sim_F w_2$ if $w_1 \sim_e w_2$ for each $e \in F$. In particular, $w_1 \sim_\emptyset w_2$ for any two epistemic worlds $w_1, w_2 \in W$.

Definition 3 *For any formula $\varphi \in \Phi$ and any epistemic world $w \in W$ of an evidence model $\langle W, E, \{\sim_e\}, \pi \rangle$, let the satisfaction relation $w \Vdash \varphi$ be defined as follows,*

1. $w \Vdash p$, if $w \in \pi(p)$,
2. $w \Vdash \neg\varphi$, if $w \not\Vdash \varphi$,
3. $w \Vdash \varphi \rightarrow \psi$, if $w \not\Vdash \varphi$ or $w \Vdash \psi$,
4. $w \Vdash A\varphi$, if there is a finite $F \subseteq E$ such that $w \sim_F u$ implies $u \Vdash \varphi$ for each $u \in W$,
5. $w \Vdash O\varphi$, if $u \Vdash \varphi$ for each $u \in W$ such that $w \sim_E u$.

3 Axioms

In addition to propositional tautologies in language Φ , our system consists of the following axioms:

1. Truth: $O\varphi \rightarrow \varphi$,

2. Positive Introspection: $A\phi \rightarrow AA\phi$,
3. Negative Introspection: $\neg O\phi \rightarrow O\neg O\phi$,
4. Distributivity: $A(\phi \rightarrow \psi) \rightarrow (A\phi \rightarrow A\psi)$ and $O(\phi \rightarrow \psi) \rightarrow (O\phi \rightarrow O\psi)$,
5. Monotonicity: $A\phi \rightarrow O\phi$.

We say that a formula ϕ is a theorem in our logical system and write $\vdash \phi$ if formula ϕ is derivable from the axioms of our systems using the Modus Ponens and the Necessitation inference rules:

$$\frac{\phi, \phi \rightarrow \psi}{\psi} \quad \frac{\phi}{A\phi}.$$

We write $X \vdash \phi$ if formula ϕ is derivable from the theorems of our logical systems and an additional set of axioms X using *only* the Modus Ponens inference rule. We say that set $X \subseteq \Phi$ is inconsistent if there is a formula $\phi \in \Phi$ such that $X \vdash \phi$ and $X \vdash \neg\phi$.

Lemma 1 *If X is a finite consistent set of formulae, and p is a propositional variable that does not occur in X , then sets $X \cup \{p\}$ and $X \cup \{\neg p\}$ are both consistent.*

PROOF. Suppose that either set $X \cup \{p\}$ or set $X \cup \{\neg p\}$ is inconsistent. Thus, either $X \vdash \neg p$ or $X \vdash p$.

First, we show that $X \vdash \neg p$ in both cases. Indeed, suppose that $X \vdash p$. Let ϕ_1, \dots, ϕ_n be the proof of formula p from the set of the assumption X . Let ϕ'_i be the result of the replacement of propositional variable p by formula $\neg p$ in formula ϕ_i . Note that variable p does not occur in set X by the assumption of the lemma. Then, ϕ'_1, \dots, ϕ'_n is the proof of formula $\neg p$ from the set of the assumption X . Therefore, $X \vdash \neg p$.

Next, we show that $X \vdash \neg\neg p$. Let ψ_1, \dots, ψ_m be the proof of formula $\neg p$ from the set of the assumption X . Consider now sequence ψ'_1, \dots, ψ'_m obtained by replacing each occurrence of variable p with $\neg p$ in proof ψ_1, \dots, ψ_m . Note again that variable p does not occur in set X by the assumption of the lemma. Then, sequence ψ'_1, \dots, ψ'_m is a proof of formula $\neg\neg p$ from the set of the assumption X . Therefore, $X \vdash \neg\neg p$.

Statements $X \vdash \neg p$ and $X \vdash \neg\neg p$ imply that set X is not consistent. \(\square\)

Lemma 2 $\vdash O\phi \rightarrow OO\phi$.

PROOF. Note that formula $O\neg O\phi \rightarrow \neg O\phi$ is an instance of the Truth axiom. Thus, $\vdash O\phi \rightarrow \neg O\neg O\phi$ by the law of contrapositive in the propositional logic. Hence, taking into account the following instance of the Negative Introspection axiom $\neg O\neg O\phi \rightarrow O\neg O\neg O\phi$, one can conclude that

$$\vdash O\phi \rightarrow O\neg O\neg O\phi. \quad (3)$$

At the same time, $\neg O\phi \rightarrow O\neg O\phi$ is an instance of the Negative Introspection axiom. Thus, $\vdash \neg O\neg O\phi \rightarrow O\phi$ by contraposition. Hence, by the Necessitation inference rule, $\vdash A(\neg O\neg O\phi \rightarrow O\phi)$. Thus, $\vdash O(\neg O\neg O\phi \rightarrow O\phi)$ by the Monotonicity axiom and the Modus Ponens inference rule. Thus, by the Distributivity axiom and the Modus Ponens inference rule, $\vdash O\neg O\neg O\phi \rightarrow OO\phi$. The last statement, together with statement (3), implies the statement of the lemma by the laws of propositional reasoning. \(\square\)

We conclude this section with the proof of the Mixed Negative Introspection axiom mentioned in Section 1.4 of the introduction.

Lemma 3 $\vdash \neg A\varphi \rightarrow O\neg A\varphi$.

PROOF. By the Truth axiom, $\vdash OA\varphi \rightarrow A\varphi$. Thus, by the law of contrapositive in the propositional logic, $\vdash \neg A\varphi \rightarrow \neg OA\varphi$. At the same time, by the Negative Introspection axiom, $\vdash \neg OA\varphi \rightarrow O\neg OA\varphi$. Hence, by the laws of propositional reasoning,

$$\vdash \neg A\varphi \rightarrow O\neg OA\varphi. \quad (4)$$

Note that $\vdash A\varphi \rightarrow AA\varphi$ by the Positive Introspection axiom and $\vdash AA\varphi \rightarrow OA\varphi$ by the Monotonicity axiom. Thus, by the laws of propositional reasoning, $\vdash A\varphi \rightarrow OA\varphi$. Hence, by the law of contrapositive in the propositional logic, $\vdash \neg OA\varphi \rightarrow \neg A\varphi$. Then, $\vdash A(\neg OA\varphi \rightarrow \neg A\varphi)$ by the Necessitation inference rule. Thus, $\vdash O(\neg OA\varphi \rightarrow \neg A\varphi)$ by the Monotonicity axiom and the Modus Ponens inference rule. Hence, $\vdash O\neg OA\varphi \rightarrow O\neg A\varphi$ by the Distributivity axiom and the Modus Ponens inference rule. Therefore, $\vdash \neg A\varphi \rightarrow O\neg A\varphi$, by the laws of propositional reasoning taking into account statement (4). \square

4 Soundness

Theorem 1 (strong soundness) *For any world $w \in W$ of each evidence model $\langle W, E, \{\sim_e\}_{e \in E}, \pi \rangle$, any set of formulae $X \subseteq \Phi$, and any formula $\varphi \in \Phi$, if $w \Vdash \chi$ for each formula $\chi \in X$ and $X \vdash \varphi$, then $w \Vdash \varphi$.*

The soundness of propositional tautologies and Modus Ponens inference rule is straightforward. Below we prove the soundness of each of the remaining axioms and the Necessitation inference rule as separate lemmas.

Lemma 4 *If $w \Vdash O\varphi$, then $w \Vdash \varphi$.*

PROOF. By Definition 3, assumption $w \Vdash O\varphi$ implies that $u \Vdash \varphi$ for all $u \in W$ such that $w \sim_E u$. Note that $w \sim_e w$ for all $e \in E$ because \sim_e is an equivalence relation. Hence, $w \sim_E w$. Therefore, $w \Vdash \varphi$. \square

Lemma 5 *If $w \Vdash A\varphi$, then $w \Vdash AA\varphi$.*

PROOF. By Definition 3, the assumption $w \Vdash A\varphi$ implies that there is a finite set $F \subseteq E$ such that $u \Vdash \varphi$ for each $u \in W$ where $w \sim_F u$.

Again by Definition 3, it suffices to show that $v \Vdash A\varphi$ for all $v \in W$ such that $w \sim_F v$. To establish this, it is enough to prove that $u \Vdash \varphi$ for all $u \in W$ such that $v \sim_F u$. Note that $w \sim_F v \sim_F u$. Thus, $w \sim_F u$ because \sim_f is an equivalence relation for each element $f \in F$. Then, $u \Vdash \varphi$ by the choice of set F . \square

Lemma 6 *If $w \Vdash \neg O\varphi$, then $w \Vdash O\neg O\varphi$.*

PROOF. By Definition 3, assumption $w \Vdash \neg O\varphi$ implies that there is $u \in W$ such that $w \sim_E u$ and $u \not\Vdash \varphi$.

Consider any $v \in W$ such that $w \sim_E v$. By Definition 3, to prove $w \Vdash O\neg O\varphi$, it suffices to show that $v \Vdash \neg O\varphi$. Note that $w \sim_E u$ and $w \sim_E v$. Thus, $v \sim_E u$ due to \sim_e being an equivalence relation for each $e \in E$. Recall that $u \not\Vdash \varphi$. Hence, $v \not\Vdash O\varphi$ by Definition 3. Therefore, $v \Vdash \neg O\varphi$ again by Definition 3. \square

Lemma 7 *If $w \Vdash A(\varphi \rightarrow \psi)$ and $w \Vdash A\varphi$, then $w \Vdash A\psi$.*

PROOF. By Definition 3, the assumption $w \Vdash A(\varphi \rightarrow \psi)$ implies that there is a finite set $F_1 \subseteq E$ such that $u \Vdash \varphi \rightarrow \psi$ for each $u \in W$ such that $w \sim_{F_1} u$. Similarly, the assumption $w \Vdash A\varphi$ implies that there is a finite set $F_2 \subseteq E$ such that $u \Vdash \varphi$ for each $u \in W$ such that $w \sim_{F_2} u$.

Let $F = F_1 \cup F_2$. It suffices to show that $u \Vdash \psi$ for each $u \in W$ such that $w \sim_F u$. Indeed, statement $w \sim_F u$ implies that $w \sim_{F_1} u$ and $w \sim_{F_2} u$. Hence, $u \Vdash \varphi \rightarrow \psi$ and $u \Vdash \varphi$ due to the choice of sets F_1 and F_2 . Therefore, $u \Vdash \psi$ by Definition 3. \square

Lemma 8 *If $w \Vdash O(\varphi \rightarrow \psi)$ and $w \Vdash O\varphi$, then $w \Vdash O\psi$.*

PROOF. Consider any $u \in W$ such that $w \sim_E u$. By Definition 3, it suffices to show that $u \Vdash \psi$. Indeed, by Definition 3, the assumptions $w \Vdash O(\varphi \rightarrow \psi)$ and $w \Vdash O\varphi$ imply that $u \Vdash \varphi \rightarrow \psi$ and $u \Vdash \varphi$. Therefore, $u \Vdash \psi$, again by Definition 3. \square

Lemma 9 *If $w \Vdash A\varphi$, then $w \Vdash O\varphi$.*

PROOF. Consider any $u \in W$ such that $w \sim_E u$. By Definition 3, it suffices to prove that $u \Vdash \varphi$. By the same definition, the assumption $w \Vdash A\varphi$ implies that there is a finite set $F \subseteq E$ such that $v \Vdash \varphi$ for each $v \in W$ such that $w \sim_F v$. Note that statement $w \sim_E u$ implies that $w \sim_F u$. Therefore, $u \Vdash \varphi$. \square

Lemma 10 *If $w \Vdash \varphi$ for each epistemic world $w \in W$ of each evidence model $\langle W, E, \{\sim_e\}_{e \in E}, \pi \rangle$, then $w \Vdash A\varphi$ for each epistemic world $w \in W$ of each evidence model $\langle W, E, \{\sim_e\}_{e \in E}, \pi \rangle$.*

PROOF. Consider any epistemic world $w \in W$ of an arbitrary evidence model $\langle W, E, \{\sim_e\}_{e \in E}, \pi \rangle$. By Definition 3, it suffices to show that there is $F \subseteq E$ such that $u \Vdash \varphi$ for each $u \in W$ where $w \sim_F u$. Indeed, let $F = \emptyset$. Note that $u \Vdash \varphi$ for each $u \in W$ due to the assumption of the lemma. \square

This concludes the proof of the soundness of our logical system.

5 Canonical Model

Theorem 2 (strong completeness) *For any set of formulae $X \subseteq \Phi$ and any formula $\varphi \in \Phi$, if $X \not\vdash \varphi$, then there is an evidence model $\langle W, E, \{\sim_e\}_{e \in E}, \pi \rangle$ and an epistemic world $w \in W$ such that $w \Vdash \chi$ for each $\chi \in X$ and $w \not\vdash \varphi$.*

As usual, the proof of a completeness theorem is using a canonical model construction. We define the canonical model in this section and give the full proof of the completeness in the appendix.

The standard proof of completeness for modal logics S4 and S5 defines worlds of a canonical model as maximal consistent sets of formulae. In our case, to specify a canonical model we need to define not only the set of worlds, but also the evidence set. In the construction that we propose, these two sets are defined to be the same set W_∞ , which is a set of nodes in a certain infinite forest. As a result, the proof of the completeness is significantly more involved than the completeness proofs for logics S4 and S5.

Throughout the section we use two operations on sequences. If w is a sequence (x_1, x_2, \dots, x_n) and u is a sequence (y_1, y_2, \dots, y_m) , then by concatenation $w :: u$ of these two sequences we mean sequence $(x_1, x_2, \dots, x_n, y_1, y_2, \dots, y_m)$. By head $hd(w)$ of a nonempty sequence $w = (x_1, x_2, \dots, x_n)$ we mean element x_n . For example, $(a, b) :: (c) = (a, b, c)$ and $hd(a, b, c) = c$.

We now define the “canonical” evidence model $\langle W_\infty, W_\infty, \{\sim_e\}_{e \in W_\infty}, \pi \rangle$. The set of epistemic worlds in the canonical model is identical to the evidence set. One can think that all pieces of evidence related in some sense to a given epistemic world are combined together into a single evidence associated with this world. Because of this, the evidence associated with world w is simply referred to as evidence w .

Informally, set W_∞ consists of sequences of the form $(X_0, \varepsilon_1, X_1, \dots, \varepsilon_n, X_n)$ where X_0, \dots, X_n are maximal consistent sets of formulae and each of $\varepsilon_1, \dots, \varepsilon_n$ is either a finite subset of W_∞ or symbol $*$. In what follows, the case when ε is a finite subset of W_∞ will form an A-accessibility relation and the case $\varepsilon = *$ will form both A-accessibility and O-accessibility relations. Such sequences can be visualised, see Figure 1, as paths in an infinite collection of infinite trees whose vertices are maximal consistent sets of formulae and whose edges are labeled with the ε 's described above.

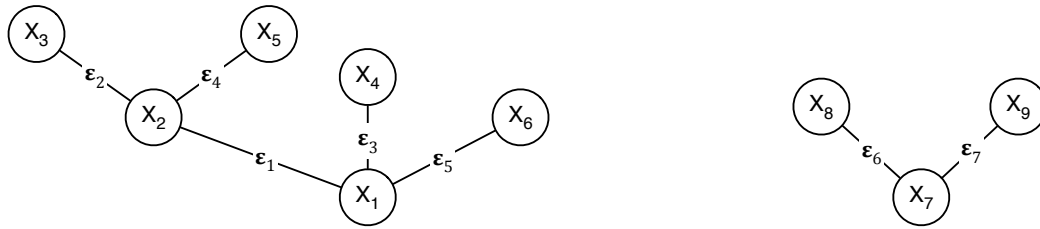


Figure 1: Fragments of the trees of sequences.

Formally, set W_∞ is specified as the union of a recursively defined infinite sequence of sets W_0, W_1, \dots . These sets represent different stages of building the infinite forest of infinite trees partially depicted in Figure 1. Note, however, that stages do not correspond to the levels of the trees. Generally speaking, vertices at the same level are not created at the same stage. After a new epistemic world is added, this world can be used as a part of evidence later in the construction.

Definition 4 A sequence of sets W_0, W_1, \dots is defined recursively as follows.

1. W_0 is the set of all single-element sequences (X_0) , where X_0 is a maximal consistent subset of Φ ,
2. W_{n+1} contains all sequences of the form $w :: (\varepsilon, Y)$ such that
 - (a) $w \in \bigcup_{i \leq n} W_i$,
 - (b) ε is either symbol $*$ or a finite subset of $\bigcup_{i \leq n} W_i$,
 - (c) Y is a maximal consistent subset of Φ ,
 - (d) if $\varepsilon = *$, then $\{\varphi \mid O\varphi \in hd(w)\} \subseteq Y$,
 - (e) if $\varepsilon \subseteq \bigcup_{i \leq n} W_i$, then $\{\varphi \mid A\varphi \in hd(w)\} \subseteq Y$.

Definition 5 $W_\infty = \bigcup_i W_i$.

Lemma 11 Set W_∞ is infinite.

PROOF. By Theorem 1, all axioms of our logical system are satisfied, in particular, in the world of any single-world model with the empty evidence set. Thus, our logical system is consistent. Hence, the empty set of formulae is consistent. Since our language contains infinitely many propositional variables, the empty set could be extended to a maximal consistent set in an infinitely many ways by Lemma 1. Hence, set W_0 is infinite by Definition 4. Therefore, set W_∞ is infinite by Definition 5. \square

Informally, two sequences are \sim_e -equivalent if they start with the same prefix and once they deviate all subsequent ε 's either are equal to $*$ or contain element e , see Figure 2. The formal definition is below.

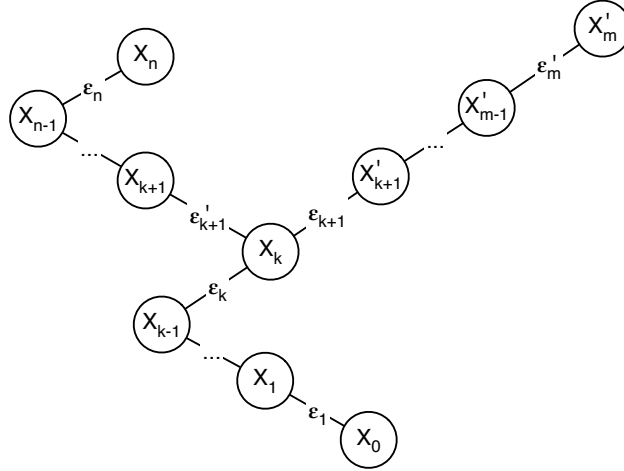


Figure 2: Illustration for Definition 6, where ε_i and ε'_i either are equal to $*$ or contain e for each $i \geq k$.

Definition 6 For any world $w = (X_0, \varepsilon_1, X_1, \dots, \varepsilon_n, X_n)$, any world $u = (X'_0, \varepsilon'_1, X'_1, \dots, \varepsilon'_m, X'_m)$, and any $e \in W_\infty$, let $w \sim_e u$ if there is k such that

1. $0 \leq k \leq \min\{n, m\}$,
2. $X_i = X'_i$ for all i such that $0 \leq i \leq k$,
3. $\varepsilon_i = \varepsilon'_i$ for all i such that $0 < i \leq k$,
4. for all i , if $k < i \leq n$, then either $e \in \varepsilon_i$ or $\varepsilon_i = *$,
5. for all i , if $k < i \leq m$, then either $e \in \varepsilon'_i$ or $\varepsilon'_i = *$.

Definition 7 $\pi(p) = \{w \in W_\infty \mid p \in hd(w)\}$.

The canonical evidence model $\langle W_\infty, W_\infty, \{\sim_e\}_{e \in W_\infty}, \pi \rangle$ is now fully defined.

6 Completeness

We start by establishing several properties of the canonical model that are necessary for our proof of the completeness. First, we show that if set $hd(w)$ contains a formula $A\varphi$, then so does set $hd(u)$ for each descendant u of vertex w .

Lemma 12 For any $k \geq 0$, any $n \geq 0$, and any $(X_0, \varepsilon_1, X_1, \dots, \varepsilon_k, X_k, \varepsilon_{k+1}, X_{k+1}, \dots, \varepsilon_{k+n}, X_{k+n}) \in W_\infty$, if $A\varphi \in X_k$, then $A\varphi \in X_{k+n}$.

PROOF. We prove the lemma by induction on n , see Figure 3 (a). If $n = 0$, then assumption $A\varphi \in X_k$ implies that $A\varphi \in X_{k+n}$. If $n > 0$, then $A\varphi \in X_{k+n-1}$ by the induction hypothesis. Hence, $X_{k+n-1} \vdash AA\varphi$ by the Positive Introspection axiom.

Case I: $\varepsilon_{k+n} = *$. Note that $X_{k+n-1} \vdash AA\varphi$ implies $X_{k+n-1} \vdash OA\varphi$ by the Monotonicity axiom. Hence, $OA\varphi \in X_{k+n-1}$ due to the maximality of set X_{k+n-1} . Then, $A\varphi \in X_{k+n}$ by Definition 4 and due to the assumption $\varepsilon_{k+n} = *$.

Case II: $\varepsilon_{k+n} \neq *$. Statement $X_{k+n-1} \vdash AA\varphi$ implies that $AA\varphi \in X_{k+n-1}$ due to the maximality of set X_{k+n-1} . Thus, $A\varphi \in X_{k+n}$ by Definition 4. \square

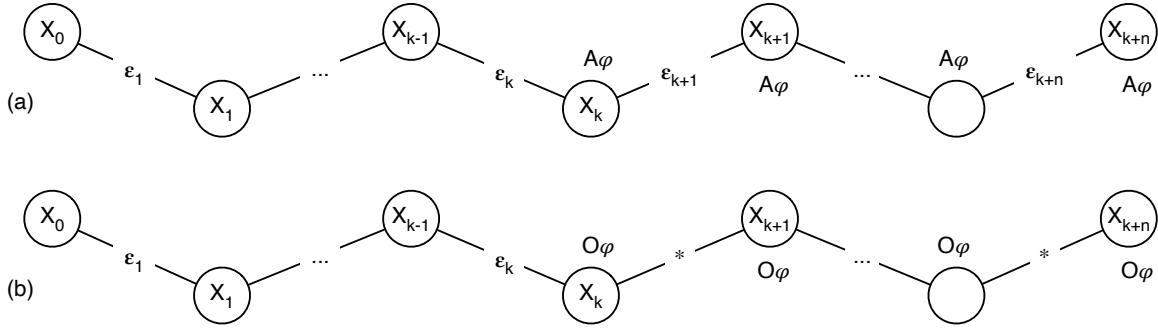


Figure 3: Illustrations: (a) for Lemma 12, and (b) for Lemma 13 and Lemma 14.

Next, we show that if $hd(w)$ contains a formula $O\phi$, then so does $hd(u)$ for each descendant u of vertex w reachable through edges all of which are labeled by symbol $*$.

Lemma 13 For any $k \geq 0$ and $n \geq 0$, and any

$$(X_0, \epsilon_1, X_1, \dots, \epsilon_k, X_k, \epsilon_{k+1}, X_{k+1}, \dots, \epsilon_{k+n}, X_{k+n}) \in W_\infty,$$

if $O\phi \in X_k$ and $\epsilon_i = *$ for all i such that $k < i \leq k+n$, then $O\phi \in X_{k+n}$.

PROOF. We prove the lemma by induction on n , see Figure 3 (b). If $n = 0$, then assumption $O\phi \in X_k$ implies that $O\phi \in X_{k+n}$. If $n > 0$, then $O\phi \in X_{k+n-1}$ by the induction hypothesis. Hence, $X_{k+n-1} \vdash OO\phi$ by Lemma 2. Thus, $O\phi \in X_{k+n}$ by Definition 4 and due to the assumption $\epsilon_{k+n} = *$. \square

The next lemma is a converse of Lemma 13. It shows that if $hd(u)$ contains a formula $O\phi$, where u is a descendant of vertex w reachable through edges all of which are labeled by symbol $*$, then so does $hd(w)$.

Lemma 14 For any $k \geq 0$, any $n \geq 0$, and any

$$(X_0, \epsilon_1, X_1, \dots, \epsilon_k, X_k, \epsilon_{k+1}, X_{k+1}, \dots, \epsilon_{k+n}, X_{k+n}) \in W_\infty,$$

if $O\phi \in X_{k+n}$ and $\epsilon_i = *$ for all i such that $k < i \leq k+n$, then $O\phi \in X_k$.

PROOF. We prove the statement of the lemma by induction on n , see again Figure 3 (b). If $n = 0$, then assumption $O\phi \in X_{k+n}$ implies that $O\phi \in X_k$. In what follows we assume that $n > 0$.

Case I: $O\phi \notin X_{k+n-1}$. Thus, $\neg O\phi \in X_{k+n-1}$ due to the maximality of the set X_{k+n-1} . Hence, $X_{k+n-1} \vdash O\neg O\phi$ by the Negative Introspection axiom. Then, $O\neg O\phi \in X_{k+n-1}$ due to the maximality of the set X_{k+n-1} . Hence, $\neg O\phi \in X_{k+n}$ by Definition 4 and because $\epsilon_{k+n} = *$. Thus, $O\phi \notin X_{k+n}$ due to the consistency of the set X_{k+n} , which is a contradiction with the assumption of the lemma.

Case II: $O\phi \in X_{k+n-1}$. Thus, $O\phi \in X_k$ by the induction hypothesis. \square

The next two lemmas are relatively standard lemmas for a completeness proof of a modal logic. Their proofs show how a sequence representing an epistemic world can be extended to produce different types of child nodes on the trees in Figure 1.

Lemma 15 For any $w \in W_\infty$, any $\neg A\phi \in hd(w)$, and any finite $F \subseteq W_\infty$, there is $u \in W_\infty$ such that $w \sim_F u$ and $\neg\phi \in hd(u)$.

PROOF. We first show that the following set is consistent: $Y_0 = \{\neg\varphi\} \cup \{\psi \mid A\psi \in hd(w)\}$. Assume the opposite. Thus, there must exist $A\psi_1, \dots, A\psi_n \in hd(w)$ such that $\psi_1, \dots, \psi_n \vdash \varphi$. Hence, by the deduction theorem for propositional logic, $\vdash \psi_1 \rightarrow (\psi_2 \rightarrow \dots (\psi_n \rightarrow \varphi) \dots)$. Then, by the Necessitation inference rule, $\vdash A(\psi_1 \rightarrow (\psi_2 \rightarrow \dots (\psi_n \rightarrow \varphi) \dots))$. By the Distributivity axiom and the Modus Ponens inference rule, $\vdash A\psi_1 \rightarrow A(\psi_2 \rightarrow \dots (\psi_n \rightarrow \varphi) \dots)$. By the Modus Ponens inference rule, $A\psi_1 \vdash A(\psi_2 \rightarrow \dots (\psi_n \rightarrow \varphi) \dots)$. By repeating the last two steps $n - 1$ times, $A\psi_1, \dots, A\psi_n \vdash A\varphi$. Hence, $hd(w) \vdash A\varphi$ by the choice of formulae ψ_1, \dots, ψ_n . Thus, $\neg A\varphi \notin hd(w)$ due to the consistency of the set $hd(w)$, which contradicts the assumption of the lemma. Therefore, set Y_0 is consistent.

Let Y be any maximal consistent extension of set Y_0 and let u be the sequence $w :: (F, Y)$. We next show that $u \in W_\infty$. Indeed, since $w \in W_\infty$, by Definition 5, there must exist $n_1 \geq 0$ such that $w \in W_{n_1}$. At the same time, since set F is a finite subset of W_∞ , by Definition 5, there must exist $n_2 \geq 0$ such that $F \subseteq \bigcup_{i \leq n_2} W_i$. Let $n = \max\{n_1, n_2\}$. Thus, $w \in \bigcup_{i \leq n} W_i$ and $F \subseteq \bigcup_{i \leq n} W_i$. Hence, $w :: (F, Y) \in W_{n+1}$ by Definition 4. Therefore, $u = w :: (F, Y) \in W_\infty$ by Definition 5. Finally, $w \sim_F u$ by Definition 6. To finish the proof of the lemma, note that $\neg\varphi \in Y_0 \subseteq Y = hd(u)$. \square

Lemma 16 For any $w \in W_\infty$ and any $\neg O\varphi \in hd(w)$, there is $u \in W_\infty$ such that $w \sim_{W_\infty} u$ and $\neg\varphi \in hd(u)$.

PROOF. We first show that the following set is consistent: $Y_0 = \{\neg\varphi\} \cup \{\psi \mid O\psi \in hd(w)\}$. Assume the opposite. Thus, there must exist formulae $O\psi_1, \dots, O\psi_n \in hd(w)$ and $\psi_1, \dots, \psi_n \vdash \varphi$. Hence, by the deduction theorem for the propositional logic, $\vdash \psi_1 \rightarrow (\psi_2 \rightarrow \dots (\psi_n \rightarrow \varphi) \dots)$. Then, by the Necessitation inference rule, $\vdash A(\psi_1 \rightarrow (\psi_2 \rightarrow \dots (\psi_n \rightarrow \varphi) \dots))$. By the Monotonicity axiom and the Modus Ponens Inference rule, $\vdash O(\psi_1 \rightarrow (\psi_2 \rightarrow \dots (\psi_n \rightarrow \varphi) \dots))$. By the Distributivity axiom and the Modus Ponens inference rule, $\vdash O\psi_1 \rightarrow O(\psi_2 \rightarrow \dots (\psi_n \rightarrow \varphi) \dots)$. By the Modus Ponens inference rule, $O\psi_1 \vdash O(\psi_2 \rightarrow \dots (\psi_n \rightarrow \varphi) \dots)$. By repeating the last two steps $n - 1$ times, $O\psi_1, \dots, O\psi_n \vdash O\varphi$. Hence, $hd(w) \vdash O\varphi$ by the choice of formulae ψ_1, \dots, ψ_n . Thus, $\neg O\varphi \notin hd(w)$ due to the consistency of the set $hd(w)$, which contradicts the assumption of the lemma. Therefore, set Y_0 is consistent.

Let Y be any maximal consistent extension of set Y_0 and let u be sequence $w :: (*, Y)$. We next show that $u \in W_\infty$. Indeed, since $w \in W_\infty$, by Definition 5, there must exist $n \geq 0$ such that $w \in W_n$. Hence, $w :: (*, Y) \in W_{n+1}$ by Definition 4. Therefore, $u = w :: (*, Y) \in W_\infty$ by Definition 5.

Finally, let us observe that $w \sim_e u$ for each $e \in W_\infty$ by Definition 6. Thus, $w \sim_{W_\infty} u$. To finish the proof of the lemma, note that $\neg\varphi \in Y_0 \subseteq Y = hd(u)$. \square

By Definition 4, elements of set W_∞ are sequences of the form $(X_0, \varepsilon_1, X_1, \dots, \varepsilon_k, X_k)$, where ε_i is either symbol $*$ or a subset of W_∞ . One might wonder if elements of W_∞ are wellfounded. In other words, is it possible for an element $w \in W_\infty$ to be a member of one of its own ε_i ? Lemma 18 shows that such elements do not exist. This is a very important observation for our proof of completeness. Lemma 17 is essentially a different form of Lemma 18 which is easier to prove by induction.

Lemma 17 For any $k \geq 0$, any $n \geq 0$, any $t \geq 0$, and any $w = (X_0, \varepsilon_1, X_1, \dots, \varepsilon_k, X_k, \varepsilon_{k+1}, X_{k+1}, \dots, \varepsilon_{k+n}, X_{k+n}) \in W_t$, if $\varepsilon_k \neq *$, then $\varepsilon_k \subseteq \bigcup_{i=0}^{t-1} W_i$.

PROOF. We prove this statement by induction on n . First, let $n = 0$. Thus, $w = (X_0, \varepsilon_1, X_1, \dots, \varepsilon_k, X_k)$. Hence, By Definition 4, assumptions $w \in W_t$ and $\varepsilon_k \neq *$ imply that $\varepsilon_k \subseteq \bigcup_{i=0}^{t-1} W_i$.

Suppose now that $n > 0$. By Definition 4, the assumption $w \in W_t$ implies that

$$(X_0, \varepsilon_1, X_1, \dots, \varepsilon_k, X_k, \varepsilon_{k+1}, X_{k+1}, \dots, \varepsilon_{k+n-1}, X_{k+n-1}) \in W_{t-1}.$$

Thus, by the induction hypothesis, $\varepsilon_k \subseteq \bigcup_{i=0}^{t-2} W_i$. Therefore, $\varepsilon_k \subseteq \bigcup_{i=0}^{t-1} W_i$. \square

Lemma 18 $w \notin \varepsilon_k$ for each $w = (X_0, \varepsilon_1, X_1, \dots, \varepsilon_n, X_n) \in W_\infty$ and each $k \leq n$.

PROOF. By Definition 5, assumption $w \in W_\infty$ implies that there is $t \geq 0$ such that $w \in W_t$. Let m be the smallest m such that $w \in W_m$. Thus, $w \notin \bigcup_{i=0}^{m-1} W_i$. At the same time, $\varepsilon_k \subseteq \bigcup_{i=0}^{m-1} W_i$ by Lemma 17. Therefore, $w \notin \varepsilon_k$. \square

The next lemma puts together the pieces of the proof that we have developed. It connects the satisfaction of a formula in an epistemic world of the canonical evidence model with the maximal consistent sets out of which the world is constructed.

Lemma 19 $w \Vdash \varphi$ iff $\varphi \in hd(w)$, for each epistemic world $w \in W_\infty$ and each formula $\varphi \in \Phi$.

PROOF. We prove the lemma by induction on the structural complexity of formula φ . If formula φ is an atomic proposition, then the required follows from Definition 7 and Definition 3. If formula φ is a negation or an implication, then the required follows from Definition 3 and the maximality and the consistency of the set $hd(w)$ in the standard way.

Suppose now that formula φ has the form $A\psi$.

(\Rightarrow) If $A\psi \notin hd(w)$, then $\neg A\psi \in hd(w)$ due to the maximality of the set $hd(w)$. To prove $w \not\Vdash A\psi$, by Definition 3, we need to show that for any finite set $F \subseteq W_\infty$ there is $u \in W_\infty$ such that $w \sim_F u$ and $u \not\Vdash \psi$. Indeed, by Lemma 15, there is $u \in W_\infty$ such that $w \sim_F u$ and $\neg\psi \in hd(u)$. Thus, $\psi \notin hd(u)$ due to the consistency of the set $hd(u)$. Therefore, $u \not\Vdash \psi$ by the induction hypothesis.

(\Leftarrow) Suppose that $A\psi \in hd(w)$. Thus, $hd(w) \vdash AA\psi$ by the Positive Introspection axiom. Hence, $hd(w) \vdash OA\psi$ by the Monotonicity axiom. Then,

$$OA\psi \in hd(w) \tag{5}$$

due to the maximality of the set $hd(w)$.

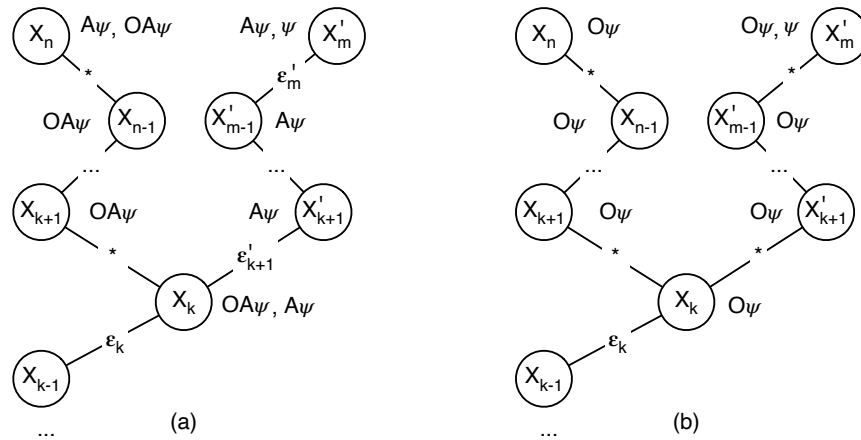


Figure 4: Illustrations for the proof of Lemma 19.

By Definition 3, it suffices to show that $u \Vdash \psi$ for all $u \in W_\infty$ such that $w \sim_w u$. By Definition 6, assumption $w \sim_w u$ implies that if $w = (X_0, \varepsilon_1, X_1, \dots, \varepsilon_n, X_n)$ and $u = (X'_0, \varepsilon'_1, X'_1, \dots, \varepsilon'_m, X'_m)$, then there is k such that,

1. $0 \leq k \leq \min\{n, m\}$,
2. $X_i = X'_i$ for all i such that $0 \leq i \leq k$,

3. $\varepsilon_i = \varepsilon'_i$ for all i such that $0 < i \leq k$,
4. for all i , if $k < i \leq n$, then either $w \in \varepsilon_i$ or $\varepsilon_i = *$,
5. for all i , if $k < i \leq m$, then either $w \in \varepsilon'_i$ or $\varepsilon'_i = *$.

By Lemma 18, $w \notin \varepsilon_i$ for each $i \leq n$. Thus, from the condition 4 above, $\varepsilon_i = *$ for each i such $k < i \leq n$, see Figure 4 (a). Hence, by Lemma 14, statement (5) implies that $\text{OA}\psi \in X_k$. Thus, $X_k \vdash \text{A}\psi$ by the Truth axiom. Then, $\text{A}\psi \in X_k$ due to the maximality of set X_k . Therefore, $\psi \in X'_m = \text{hd}(u)$ by Lemma 12 and condition 5 above. Therefore, $u \Vdash \psi$ by the induction hypothesis.

Finally, let formula φ have the form $\text{O}\psi$.

(\Rightarrow) If $\text{O}\psi \notin \text{hd}(w)$, then $\neg\text{O}\psi \in \text{hd}(w)$ due to the maximality of the set $\text{hd}(w)$. To prove $w \not\Vdash \text{O}\psi$, by Definition 3, we need to show that $u \not\Vdash \psi$ for some $u \in W_\infty$ such that $w \sim_{W_\infty} u$. Indeed, by Lemma 16, there is $u \in W_\infty$ such that $w \sim_{W_\infty} u$ and $\neg\psi \in \text{hd}(u)$. Thus, $\psi \notin \text{hd}(u)$ due to the consistency of the set $\text{hd}(u)$. Therefore, $u \not\Vdash \psi$ by the induction hypothesis.

(\Leftarrow) Suppose that $\text{O}\psi \in \text{hd}(w)$. By Definition 3, it suffices to show that $u \Vdash \psi$ for all $u \in W_\infty$ such that $w \sim_{W_\infty} u$. Indeed, let $w = (X_0, \varepsilon_1, X_1, \dots, \varepsilon_n, X_n)$ and $u = (X'_0, \varepsilon'_1, X'_1, \dots, \varepsilon'_m, X'_m)$. By Definition 6, assumption $w \sim_{W_\infty} u$ implies that for each $e \in W_\infty$ there is integer k_e such that

1. $0 \leq k_e \leq \min\{n, m\}$,
2. $X_i = X'_i$ for all i such that $0 \leq i \leq k_e$,
3. $\varepsilon_i = \varepsilon'_i$ for all i such that $0 < i \leq k_e$,
4. for all i , if $k_e < i \leq n$, then either $w \in \varepsilon_i$ or $\varepsilon_i = *$,
5. for all i , if $k_e < i \leq m$, then either $w \in \varepsilon'_i$ or $\varepsilon'_i = *$.

By Lemma 11, set W_∞ is infinite and, thus, it is nonempty. Hence, set $\{k_e \mid e \in W_\infty\}$ is nonempty. Set $\{k_e \mid e \in W_\infty\}$ is finite because $0 \leq k_e \leq \min\{n, m\}$ for each $e \in W_\infty$. Let k be the maximal element of the set $\{k_e \mid e \in W_\infty\}$. Hence,

1. $0 \leq k \leq \min\{n, m\}$,
2. $X_i = X'_i$ for all i such that $0 \leq i \leq k$,
3. $\varepsilon_i = \varepsilon'_i$ for all i such that $0 < i \leq k$,
4. for all i , if $k < i \leq n$, then either $W_\infty \subseteq \varepsilon_i$ or $\varepsilon_i = *$,
5. for all i , if $k < i \leq m$, then either $W_\infty \subseteq \varepsilon'_i$ or $\varepsilon'_i = *$.

By Lemma 11, set W_∞ is infinite yet sets ε'_i and ε_i are finite for each i by Definition 4. Hence, $W_\infty \not\subseteq \varepsilon_i$ and $W_\infty \not\subseteq \varepsilon'_i$ for each i . Thus, conditions 4 and 5 above imply that $\varepsilon_i = *$ for all i such that $k < i \leq n$ and $\varepsilon'_i = *$ for all i such that $k < i \leq m$, see Figure 4 (b). Thus, by Lemma 14, assumption $\text{O}\psi \in \text{hd}(w) = X_n$ implies that $\text{O}\psi \in X_k = X'_k$. Therefore, $\psi \in X'_m = \text{hd}(u)$, by Lemma 13. \square

We are now ready to finish the proof of Theorem 2. Set $X \cup \{\neg\varphi\}$ is consistent by the assumption $X \not\vdash \varphi$ of the theorem. Consider any maximal consistent extension X_0 of set $X \cup \{\neg\varphi\}$. Let w_0 be the single-element sequence (X_0) . Thus, $w_0 \Vdash \chi$ for each formula $\chi \in X$ and $w_0 \not\Vdash \varphi$ by Lemma 19. This concludes the proof of the theorem.

7 Conclusion

In this paper we have shown that knowledge obtained from infinitely many pieces of evidence has properties captured by modal logic S5 and knowledge obtained from finitely many pieces of evidence has properties described by modal logic S4. The main technical result is a sound and complete propositional bi-modal logic that captures properties of both of these types of knowledge and their interplay.

A natural next step is to consider first-order logic with the same two modalities. Note that the knowledge modality O satisfies Barcan Formula [3] $\forall x O\varphi \rightarrow O\forall x\varphi$ because this formula is derivable from S5 axioms stated in the first-order modal language [14]. At the same time, attainable knowledge modality A does not satisfy Barcan Formula $\forall x A\varphi \rightarrow A\forall x\varphi$. Indeed, if variable x ranges over an infinite domain and each value of x in this domain has a distinct single evidence e_x that justifies $\varphi(x)$, then $\forall x A\varphi$ is true, but $A\forall x\varphi$ is not. A complete axiomatization of the interplay of these two modalities in the first-order language remains an open problem.

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